

***m* -Hypergeometric Solutions of Anti-Difference and
q -Anti-Difference Equations**

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Abstract

In this paper we consider the problem of finding *m*-hypergeometric solutions of anti-difference equations. We extend the greatest factorial factorization (GFF) of a polynomial, introduced by Paule (1995), to the *m*-greatest factorial factorization (*m* GFF). Equipped with the *m* GFF-concept, we present algebraically motivated approach to the problem. This approach requires only “gcd” operations but no factorization. Then, we solve the same problem for *q*-anti-difference equations.

Keywords : Gosper’s algorithm, *m* -hypergeometric solution, m-greatest factorial factorization, *q* -Gosper algorithm, *qm* -hypergeometric solution, *qm* -greatest factorial factorization.

q -

m-
gcd
m-
(1995)
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1. Introduction

Let m denotes a positive integer, \mathbb{N} be the set of natural numbers, K be the field of characteristic zero, $K(n)$ be the field of rational functions over K , $K[n]$ be the ring of polynomials over K , F denotes the transcendental extension of K by the indeterminate q , i.e., $F = K(q)$, E denotes the shift operator on $K[n]$, i.e., $(Ep)(n) = p(n+1)$ for any $p \in K[n]$, ε denotes the q -shift operator on $F[n]$ and $F(n)$, i.e., $(\varepsilon u)(n) = u(qn)$ for any $u \in F[n]$ or $u \in F(n)$, $\deg(p)$ denotes the polynomial degree (in n) of any $p \in K[n]$ or $p \in F[n]$, $p \neq 0$. We define $\deg(0) = -1$. We assume the result of any gcd (greatest common divisor) computation in $K[n]$ or $F[n]$ as being normalized to a monic polynomial p , i.e., the leading coefficient of p being 1. Recall that a non-zero term t_n is called a hypergeometric term over K if there exist a rational function $r(n) \in K(n)$ such that

$$\frac{t_{n+1}}{t_n} = r(n).$$

Gosper's algorithm (Gosper, 1978) (also see Graham *et al.*, 1989, Koepf, 1998, Petkovšek *et al.*, 1996) has been extensively studied and widely used to prove hypergeometric identities. Given a hypergeometric term t_n , Gosper's algorithm is a procedure to find a hypergeometric term z_n satisfying

$$z_{n+1} - z_n = t_n. \quad (1.1)$$

if it exists, or confirm the nonexistence of any solution of (1.1). In Paule (1995), Paule introduced the GFF-concept. Equipped with the GFF-concept, he presented a new and algebraically motivated approach to Gosper's algorithm.

A non-zero term a_n is called an m -hypergeometric over K if there exist a rational function $w(n) \in K[n]$ such that

$$\frac{a_{n+m}}{a_n} = w(n). \quad (1.2)$$

In Koepf (1995), Koepf extends Gosper's algorithm to find m -hypergeometric solutions h_n of

$$h_{n+m} - h_n = a_n, \quad (1.3)$$

where a_n is a given m -hypergeometric term. In Petkovšek and Bruno (1993), Petkovšek and Bruno described an algorithm to find m -hypergeometric solutions of homogeneous linear recurrences with polynomial coefficients. Their algorithm reduces to algorithm **Hyper** (Petkovšek, 1992) when $m = 1$.

A non-zero term b_k is called a q -hypergeometric over F if there exists a rational function $\sigma \in F(q^k)$ such that

$$\frac{b_{k+1}}{b_k} = \sigma(q^k).$$

In Paule and Riese (1997), Paule and Riese introduced the q -greatest factorial factorization (q GFF) of polynomials, which is a q -analogue of the GFF-concept. Equipped with the q GFF, they presented a new approach to find q -hypergeometric solutions l_k of

$$l_{k+1} - l_k = b_k, \tag{1.4}$$

where b_k is a given q -hypergeometric term. Paule-Riese's approach can be viewed as an q -analogue of Gosper's algorithm.

A non-zero term f_k is called a qm -hypergeometric over F if there exist a rational function $\rho \in F(q^k)$ such that

$$\frac{f_{k+m}}{f_k} = \rho(q^k).$$

Let us define the dispersion $\text{dis}_m(a, b)$ of the polynomials $a(n), b(n) \in K[n]$ to be the greatest nonnegative integer k (if it exists) such that $a(n)$ and $b(n + mk)$ have a nontrivial common divisor, i.e.,

$$\text{dis}_m(a, b) = \max \{ k \in \mathbb{N} \mid \deg \gcd(a(n), b(n + mk)) \geq 1 \}.$$

If k does not exist then we set $\text{dis}_m(a, b) = -1$. Recall that the pair $\langle c, d \rangle, c, d \in K[n]$, is called the reduced form of $r \in K(n)$ if $r = \frac{c}{d}$, d is monic, and $\gcd(c, d) = 1$.

The contents of this paper are as follows: In Section 2, we give the Fundamental m GFF Lemma, which is an extension of the Fundamental Lemma given by Paule (1995). In Section 3, we extend Paule's approach to find m -hypergeometric solutions of anti-difference equations. In Section 4, we give the Fundamental qm GFF Lemma, which is an extension of the Fundamental q GFF Lemma given by Paule and Riese (1997). Finally, In Section 5, we extend Paule-Riese's approach to find qm -hypergeometric solutions of q -anti-difference equations.

2. m -Greatest Factorial Factorization

In this section we define the m GFF of a polynomial, which is an extension of the GFF-concept introduced by Paule.

2.1 Basic Definitions

Definition 2.1. For any monic polynomial $p \in K[n]$ and $i \in \mathbb{N}$, the i -th m -falling factorial $[p]_m^i$ of p is defined as

$$[p]_m^i = p \cdot E^{-m} p \cdot E^{-2m} p \cdot \dots \cdot E^{-(i+1)m} p.$$

For $i = 0$, we let $[p]_m^0 = 1$.

Definition 2.2. We say that $\langle p_1, p_2, \dots, p_s \rangle$, $p_i \in K[n]$, is an m GFF-form of a monic polynomial $p \in K[n]$ if the following conditions hold:

- (m GFF1) $p = [p_1]_m^1 \cdot [p_2]_m^2 \cdot \dots \cdot [p_s]_m^s$,
- (m GFF2) each p_i is monic and $s > 0$ implies $\deg(p_s) > 0$,
- (m GFF3) $\gcd([p_i]_m^i, E^m p_j) = 1$ for $1 \leq i \leq j \leq s$,
- (m GFF4) $\gcd([p_i]_m^i, E^{-jm} p_j) = 1$ for $1 \leq i \leq j \leq s$.

If $\langle p_1, p_2, \dots, p_s \rangle$ is an m GFF-form of a monic $p \in K[n]$ we sometimes express this fact for short by m GFF(p) = $\langle p_1, p_2, \dots, p_s \rangle$.

2.2 The Fundamental m GFF Lemma

In this section we give the Fundamental m GFF Lemma, which is an extension of the Fundamental Lemma given by Paule. The $\gcd(p, E^m p)$ for $p \in K[n]$ plays a basic role in finding m -hypergeometric solutions of anti-difference equation (1.3).

Lemma 2.1. (“Fundamental m GFF Lemma”) Given a monic polynomial $p \in K[n]$ with m GFF-form $\langle p_1, p_2, \dots, p_s \rangle$. Then

$$\gcd(p, E^m p) = [p_2]_m^1 \cdot [p_3]_m^2 \cdot \dots \cdot [p_s]_m^{s-1}.$$

Proof. Proceeding by induction on s the case $s = 0$ is trivial. For $s > 0$,

$$\begin{aligned} \gcd(p, E^m p) &= [p_s]_m^{s-1} \cdot \gcd([p_1]_m^1 \cdot \dots \cdot [p_{s-1}]_m^{s-1} \cdot E^{(-s+1)m} p_s, E^m ([p_1]_m^1 \cdot \dots \cdot [p_{s-1}]_m^{s-1} \cdot p_s)). \\ &= [p_s]_m^{s-1} \cdot \gcd([p_1]_m^1 \cdot \dots \cdot [p_{s-1}]_m^{s-1}, E^m ([p_1]_m^1 \cdot \dots \cdot [p_{s-1}]_m^{s-1})). \end{aligned}$$

The first equality is obvious, the second is a consequence of m GFF3 and m GFF4 because for $i < s$ we have

$$\gcd([p_i]_m^i, E^m p_s) = \gcd(E^{(-s+1)m} p_s, E^m [p_i]_m^i) = E^m \gcd(E^{-sm} p_s, [p_i]_m^i) = 1.$$

together with $\gcd(E^{(-s+1)^m} p_s, E^m p_s) \mid \gcd([p_s]_m^s, E^m p_s) = 1$. The rest of the proof follows from applying the induction hypothesis.

□

In the above lemma we see that from the m GFF-form of a polynomial p we can find the m GFF-form of $\gcd(p, E^m p)$.

3. m -Hypergeometric Solutions of Anti-Difference Equations

In this section we extend Paule approach to find m -hypergeometric solutions h_n of equation (1.3). Given an m -hypergeometric term a_n and suppose that there exists an m -hypergeometric term h_n satisfying equation (1.3), then by using (1.3) we find

$$\frac{h_n}{a_n} = \frac{h_n}{h_{n+m} - h_n} = \frac{1}{\frac{h_{n+m}}{h_n} - 1}.$$

Let $y(n) = \frac{h_n}{a_n}$. It follows that $y(n)$ is a rational function of n . Let $\langle a, b \rangle$ be the reduced form of $w(n) = \frac{a_{n+m}}{a_n}$. Substituting $y(n)a_n$ for h_n in (1.3) to obtain

$$a(n)y(n+m) - b(n)y(n) = b(n). \tag{3.1}$$

This means, the problem of finding m -hypergeometric solution of (1.3) is equivalent to finding a rational solution $y(n)$ of (3.1). If a solution $y(n) \in K(n)$ of (3.1) with the reduced form $\langle u, v \rangle$ exists, assume we know v or a multiple $V \in K[n]$ of v . Then equation (3.1) can be written as

$$a(n) \cdot V(n) \cdot U(n+m) - b(n) \cdot V(n+m) \cdot U(n) = b(n) \cdot V(n) \cdot V(n+m), \tag{3.2}$$

where $U(n) = \frac{u(n) \cdot V(n)}{v(n)}$ is unknown polynomial. Hence the problem reduces to finding a polynomial solution $U \in K[n]$ of equation (3.2). To solve (3.2) we try to find a suitable denominator polynomial V and then U can be computed as a polynomial solution of (3.2). Let

$$v_i(n) = \frac{v(n+i)}{\gcd(v, E^m v)} \quad \text{for } i \in \{0, m\}.$$

Then (3.1) is equivalent to

$$(3.3) \quad a(n) \cdot v_0(n) \cdot u(n+m) - b(n) \cdot v_m(n) \cdot u(n) = b(n) \cdot v_0(n) \cdot v_m(n) \cdot \gcd(v, E^m v).$$

From this equation we immediately get that $v_0(n) \mid b(n)$ and that $v_m(n) \mid a(n)$. Let m GFF $(v) = \langle p_1, p_2, \dots, p_s \rangle$, by using the m GFF-concept and the Fundamental m GFF Lemma we get that

$$v_0 = \frac{v}{\gcd(v, E^m v)} = p_1 \cdot E^{-m} p_2 \dots E^{(-s+1)m} p_s \mid b(n), \quad (3.4)$$

$$v_m = \frac{E^m v}{\gcd(v, E^m v)} = E^m p_1 \cdot E^m p_2 \dots E^m p_s \mid a(n), \quad (3.5)$$

This observation gives rise to a simple algorithm for computing a multiple $V = [P_1]_m^1 \cdot [P_2]_m^2 \cdots [P_s]_m^s$ of v .

- Straightforward conclusion

$$p_i \mid \gcd(E^{-m} a, E^{(i-1)m} b) \quad \forall i \in \{1, \dots, s\}.$$

If $P_i = \gcd(E^{-m} a, E^{(i-1)m} b)$ then obviously $p_i \mid P_i$. Thus, we could take

$$m \text{ GFF}(V) = \langle P_1, P_2, \dots, P_N \rangle,$$

where $N = \text{dis}_m(a, b) = \max \{ i \in \mathbb{N} \mid \deg \gcd(a, E^{im} b) \geq 1 \}$. If N is not defined then we set $V = 1$.

- Refined conclusions

$$p_1 \mid \gcd(E^{-m} a, b).$$

If $P_1 = \gcd(E^{-m} a, b)$ then $p_1 \mid P_1$ and

$$p_2 \mid \gcd\left(E^{-m} \left(\frac{a}{E^m(P_1)}\right), E^m \left(\frac{b}{P_1}\right)\right).$$

If $P_2 = \gcd\left(E^{-m} \left(\frac{a}{E^m(P_1)}\right), E^m \left(\frac{b}{P_1}\right)\right)$, then $p_{12} \mid P_2$ and so on until we arrive at a P_N and we may again take m GFF $(V) = \langle P_1, P_2, \dots, P_N \rangle$.

□

The algorithm that we have just derived for (1.3) can be written, by using the “redefined conclusions”, as follows:

Algorithm 3.1.

INPUT : $w(n) \in K(n)$ such that $\frac{a_{n+m}}{a_n} = w(n)$ for all $n \in \mathbb{N}$.

OUTPUT: an m -hypergeometric solution h_n of (1.3) if it exists, otherwise “no m -hypergeometric solution of (1.3) exists”.

(1) Decompose $w(n)$ into the reduced form $\langle a, b \rangle$.

(2) Compute $N = \text{dis}_m(a, b) = \max \{i \in \mathbb{N} \mid \deg \gcd(a, E^{im}b) \geq 1\}$.

If $N > 0$ then compute for j from 1 to N

$$P_j(n) = \gcd(E^{-m}a, E^{(j-1)m}b)$$

$$a = \frac{a}{E^m P_j(n)}$$

$$b = \frac{b}{E^{-(j-1)m} P_j(n)}$$

$$m \text{ GFF}(V) = \langle P_1, P_2, \dots, P_N \rangle$$

otherwise $V = 1$.

(3) If equation (3.2) can be solved for $U \in K[n]$ then return $h_n = \frac{U(n)}{V(n)} a_n$, otherwise return “no m -hypergeometric solution of (1.3) exists”.

□

4. q - m -Greatest Factorial Factorization

In this section we define the “ q - m -Greatest Factorial Factorization” (qm GFF) of a polynomial which is an extension of the q GFF-concept introduced by Paule and Riese. Also, it is a q -analogue of the m GFF-concept, defined with respect to the q -shift operator ε instead of the shift operator E as for m GFF. In Sections 4 and 5, we will use n as an abbreviation for q^k .

Let us define the dispersion $\text{dis}_m(a(n), b(n))$ of the q -monic polynomials $a(n), b(n) \in F[n]$ is the greatest nonnegative integer i (if it exists) such that $a(n)$ and $b(q^{mi}n)$ have a nontrivial common divisor, i.e.,

$$\text{dis}_m(a, b) = \max \{i \in \mathbb{N} \mid \deg \gcd(a(n), b(q^{mi}n)) \geq 1\}.$$

A polynomial $p \in F[n]$ is said to be q -monic if $p(0) = 1$. Any polynomial $p \in F[n]$ has a unique factorization, the q -monic decomposition, in the form

$$p = z \cdot n^\alpha \cdot \hat{p},$$

where $z \in F$, $\alpha \in \mathbb{N}$, and $\hat{p} \in F[n]$ q -monic. We will write \gcd_q instead of “gcd”, indicating that the \gcd_q of two q -monic polynomials is understood to be q -monic.

More generally, if $p_1 = z_1 \cdot n^{\alpha_1} \cdot \hat{p}_1$ and $p_2 = z_2 \cdot n^{\alpha_2} \cdot \hat{p}_2$ are q -monic decompositions of $p_1, p_2 \in F[n]$, we define

$$\gcd_q(p_1, p_2) = \gcd(n^{\alpha_1}, n^{\alpha_2}) \cdot \gcd_q(\hat{p}_1, \hat{p}_2).$$

4.1 Basic Definitions

Definition 4.1. For any q -monic polynomial $p \in F[n]$ and $i \in \mathbb{N}$, the i -th m -falling q -factorial $[p]_{m_q}^i$ of p is defined as

$$[p]_{m_q}^i = p \cdot \varepsilon^{-m} p \cdot \varepsilon^{-2m} p \cdots \varepsilon^{-(i-1)m} p.$$

For $i = 0$, we let $[p]_{m_q}^0 = 1$.

Definition 4.2. We say that $\langle p_1, p_2, \dots, p_s \rangle$, $p_i \in F[n]$, is an qm GFF-form of a q -monic polynomial $p \in F[n]$ if the following conditions hold:

- (qm GFF1) $p = [p_1]_{m_q}^1 \cdot [p_2]_{m_q}^2 \cdots [p_s]_{m_q}^s$,
- (qm GFF2) each p_i is q -monic and $s > 0$ implies $\deg(p_s) > 0$,
- (qm GFF3) $\gcd_q([p_i]_{m_q}^i, \varepsilon^m p_j) = 1$ for $1 \leq i \leq j \leq s$,
- (qm GFF4) $\gcd_q([p_i]_{m_q}^i, \varepsilon^{-jm} p_j) = 1$ for $1 \leq i \leq j \leq s$.

If $\langle p_1, p_2, \dots, p_s \rangle$ is the qm GFF-form of a q -monic $p \in F[n]$ we also denote this fact for short by qm GFF(p) = $\langle p_1, p_2, \dots, p_s \rangle$.

4.2 The Fundamental qm GFF Lemma

In this section we give the Fundamental qm GFF Lemma which is an extension of the Fundamental q GFF Lemma given by Paule and Riese. In finding m -hypergeometric solutions of anti-difference equations (i.e., $q = 1$) the $\gcd(p, E^m p)$ for $p \in K[n]$ plays a basic role. The same is true for qm -hypergeometric solutions of q -anti-difference equations with respect to the q -shift operator ε instead of E .

Lemma 4.1. (“Fundamental qm GFF Lemma”) *Given a q -monic polynomial $p \in F[n]$ with qm GFF-form $\langle p_1, p_2, \dots, p_s \rangle$. Then*

$$\gcd_q(p, \varepsilon^m p) = [p_2]_{m_q}^1 \cdot [p_3]_{m_q}^2 \cdots [p_s]_{m_q}^{s-1}.$$

Proof. Proceeding by induction on s the case $s = 0$ is trivial. For $s > 0$,

$$\begin{aligned} \gcd_q(p, \varepsilon^m p) &= [p_s]_{m_q}^{s-1} \cdot \gcd_q([p_1]_{m_q}^1 \cdots [p_{s-1}]_{m_q}^{s-1} \cdot \varepsilon^{(-s+1)m} p_s, \varepsilon^m ([p_1]_{m_q}^1 \cdots [p_{s-1}]_{m_q}^{s-1} \cdot p_s)). \\ &= [p_s]_{m_q}^{s-1} \cdot \gcd_q([p_1]_{m_q}^1 \cdots [p_{s-1}]_{m_q}^{s-1}, \varepsilon^m ([p_1]_{m_q}^1 \cdots [p_{s-1}]_{m_q}^{s-1})). \end{aligned}$$

The first equality is obvious, the second is a consequence of qm GFF3 and qm GFF4 because for $i < s$ we have

$$\gcd_q([p_i]_{m_q}^i, \varepsilon^m p_s) = \gcd_q(\varepsilon^{(-s+1)m} p_s, \varepsilon^m [p_i]_{m_q}^i) = \varepsilon^m \gcd_q(\varepsilon^{-sm} p_s, [p_i]_{m_q}^i) = 1.$$

together with $\gcd_q(\varepsilon^{(-s+1)m} p_s, \varepsilon^m p_s) \mid \gcd_q([p_s]_{m_q}^s, \varepsilon^m p_s) = 1$. The rest of the proof follows from applying the induction hypothesis.

□

In the above lemma we see that from the qm GFF-form of a q -monic polynomial p one directly can extract the qm GFF-form of $\gcd_q(p, \varepsilon^m p)$.

5. qm -Hypergeometric Solutions of q -Anti-Difference Equations

In this section we extend Paule-Riese's approach to find qm -hypergeometric solutions g_k of the q -anti-difference equation

$$g_{k+m} - g_k = f_k, \tag{5.1}$$

where f_k is a given a qm -hypergeometric term. Given a qm -hypergeometric term f_k and suppose that there exists a qm -hypergeometric term f_k satisfying equation (5.1), then by using (5.1) we find

$$\frac{g_k}{f_k} = \frac{g_k}{g_{k+m} - g_k} = \frac{1}{\frac{g_{k+m} - 1}{g_k}}.$$

Let $\tau = \frac{g_k}{f_k}$. It follows that τ is a rational function over F . Substituting $\tau \cdot f_k$ for g_k in (5.1) to obtain

$$\rho \cdot \varepsilon^m \tau - \tau = 1, \quad (5.2)$$

where $\rho = \frac{f_{k+m}}{f_k} \in F(n)$ is a rational function. Let $\rho = z \cdot n^\alpha \cdot \frac{a}{b}$ with $z \in F$, α integer, and $a, b \in F[n]$ relatively prime and q -monic. For any integer α we define $\alpha_+ = \max(\alpha, 0)$ and $\alpha_- = \max(-\alpha, 0)$, thus equation (5.2) is equivalent to

$$z \cdot n^{\alpha_+} \cdot a \cdot \varepsilon^m \tau - n^{\alpha_-} \cdot b \cdot \tau = n^{\alpha_-} \cdot b. \quad (5.3)$$

This means the problem of finding a qm -hypergeometric solutions of (5.1) is equivalent to finding rational solutions τ of (5.3). Let $\tau = \frac{u}{v}$ where $u, v \in F[n]$ be two unknown relatively prime polynomials with $v = n^\beta \cdot \hat{v}$ the q -monic decomposition of v . If a solution τ of (5.3) exists, assume we know v or a multiple $V \in F[n]$ of v . Then equation (5.3) can be written as

$$z \cdot n^{\alpha_+} \cdot a \cdot V \cdot \varepsilon^m U - n^{\alpha_-} \cdot b \cdot \varepsilon^m V \cdot U = n^{\alpha_-} \cdot b \cdot V \cdot \varepsilon^m V. \quad (5.4)$$

where $U(n) = \frac{u(n) \cdot V(n)}{v(n)}$ is unknown polynomial. Hence the problem reduces to finding a polynomial solution $U \in F[n]$ of equation (5.4). To solve (5.4) we try to find a suitable denominator polynomial V and then U can be computed as a polynomial solution of (5.4). Let

$$v_i = \frac{\varepsilon^i v}{\gcd_q(v, \varepsilon^m v)} \quad \text{for } i \in \{0, m\}.$$

Then (5.3) is equivalent to

$$z \cdot n^{\alpha_+} \cdot a \cdot v_0 \cdot \varepsilon^m u - n^{\alpha_-} \cdot b \cdot v_m \cdot u = n^{\alpha_-} \cdot b \cdot v_0 \cdot v_m \cdot \gcd_q(v, \varepsilon^m v). \quad (5.5)$$

From this equation we immediately get that $v_0 \mid b$ and that $v_m \mid a$. Let qm GFF $(\hat{v}) = \langle p_1, p_2, \dots, p_s \rangle$, by using the qm GFF-concept and the Fundamental qm GFF Lemma we get that

$$v_0 = \frac{v}{\gcd_q(v, \varepsilon^m v)} = \frac{\hat{v}}{\gcd_q(\hat{v}, \varepsilon^m \hat{v})} = p_1 \cdot \varepsilon^{-m} p_2 \dots \varepsilon^{(-s+1)m} p_s \mid b, \quad (5.6)$$

$$v_m = \frac{\varepsilon^m v}{\gcd_q(v, \varepsilon^m v)} = \frac{q^{m\beta} \cdot \varepsilon^m \hat{v}}{\gcd_q(\hat{v}, \varepsilon^m \hat{v})} = q^{m\beta} \cdot \varepsilon^m p_1 \cdot \varepsilon^m p_2 \dots \varepsilon^m p_s \mid a, \tag{5.7}$$

This observation give rise to a simple algorithm for computing a multiple $\hat{V} = [P_1]_{m_q}^1 \cdot [P_2]_{m_q}^2 \dots [P_s]_{m_q}^s$ of \hat{v} .

- Straightforward conclusion

$$p_i \mid \gcd_q(\varepsilon^{-m} a, \varepsilon^{(i-1)m} b) \quad \forall i \in \{1, \dots, s\}.$$

If $P_i = \gcd_q(\varepsilon^{-m} a, \varepsilon^{(i-1)m} b)$ then obviously $p_i \mid P_i$. Thus, we could take

$$\hat{V} = [P_1]_{m_q}^1 \cdot [P_2]_{m_q}^2 \dots [P_N]_{m_q}^N,$$

where $N = \text{dis}_m(a, b) = \max \{i \in \mathbb{N} \mid \deg \gcd_q(a, \varepsilon^{im} b) \geq 1\}$. If N is not defined then we set $\hat{V} = 1$.

- Refined conclusions

$$p_1 \mid \gcd_q(\varepsilon^{-m} a, b).$$

if we take $P_1 = \gcd_q(\varepsilon^{-m} a, b)$ then $p_1 \mid P_1$ and

$$p_2 \mid \gcd_q\left(\varepsilon^{-m} \left(\frac{a}{\varepsilon^m(P_1)}\right), \varepsilon^m \left(\frac{b}{P_1}\right)\right).$$

If we take $P_2 = \gcd_q\left(\varepsilon^{-m} \left(\frac{a}{\varepsilon^m(P_1)}\right), \varepsilon^m \left(\frac{b}{P_1}\right)\right)$, then $p_2 \mid P_2$ and so on until

we arrive at a P_N and we may again take $\hat{V} = [P_1]_{m_q}^1 \cdot [P_2]_{m_q}^2 \dots [P_N]_{m_q}^N$.

□

With \hat{V} in hand, all what is left for solving (5.4), and thus finding the a qm -hypergeometric solution of equation (5.1), is to determine an appropriate value of γ such that

$$v(n) = n^\beta \cdot \hat{v}(n) \mid V(n) = n^\gamma \cdot \hat{V}(n).$$

For that we will follow the approach given by Paule and Riese (1997). Consider equation (5.5): (i) Assume that $\alpha \neq 0$ then either $\alpha_- \neq 0$ or $\alpha_+ \neq 0$. In the first case we have $\alpha_+ = 0$ and $n^{\alpha_-} \mid u$, hence β must be zero because of $\gcd_q(u, v) = 1$. This

means we can choose $\gamma = 0$. In the second case we have $\alpha_- = 0$ and $n^{\min(\alpha_+, \beta)} \mid u$, because of $\gcd_q(v, \varepsilon^m v) = n^\beta \cdot \gcd_q(\hat{v}, \varepsilon^m \hat{v})$. Again β must be zero, and again we can choose $\gamma = 0$. (ii) Assume that $\alpha = 0$. In this case equation (5.5) evaluated at $n = 0$ turns into

$$(z - q^{\beta m})u(0) = q^{\beta m} \cdot \delta_{0, \beta},$$

where $\delta_{0, \beta}$ denotes the Kronecker symbol. This means if $\beta > 0$ we obtain, observing that $u(0) \neq 0$ in this case, as a condition for β that $z = q^{\beta m}$. Hence in case $\alpha = 0$, we choose $\gamma = \frac{1}{m} \cdot \log_q(z)$ if z is a positive integer power of q , or $\gamma = 0$ otherwise.

The algorithm that we have just derived for (5.1) can be written, by using the “redefined conclusions”, as follows:

Algorithm 5.1.

INPUT : $\rho \in F(n)$ such that $f_{k+m}/f_k = \rho(q^k)$ for all $n \in \mathbb{N}$.

OUTPUT: a *qm*-hypergeometric solution g_k of (5.1) if it exists, otherwise “no *qm*-hypergeometric solution g_k of (5.1) exists”.

(1) Decompose ρ into the form $\rho = z \cdot n^\alpha \cdot \frac{a}{b}$ such that $z \in F$, α integer, and $a, b \in F[n]$ relatively prime and q -monic.

(2) Compute $N = \text{dis}_m(a, b) = \max \{i \in \mathbb{N} \mid \deg \gcd_q(a, \varepsilon^{im} b) \geq 1\}$.

If $N > 0$ then compute for j from 1 to N

$$P_j(n) = \gcd_q(\varepsilon^{-m} a, \varepsilon^{(j-1)m} b)$$

$$a = \frac{a}{\varepsilon^m P_j(n)}$$

$$b = \frac{b}{\varepsilon^{-(j-1)m} P_j(n)}$$

$$\hat{V} = [P_1]_{m_q}^1 \cdot [P_2]_{m_q}^2 \cdots [P_N]_{m_q}^N$$

otherwise $\hat{V} = 1$.

(3) Determine the value of γ as follows:

$$\gamma = \begin{cases} \frac{1}{m} \cdot \log_q(z) & \text{if } \alpha = 0 \text{ and } z \text{ a positive integer power of } q, \\ 0 & \text{otherwise.} \end{cases}$$

(4) Take $V = n^\gamma \cdot \hat{V}$. If equation (5.4) can be solved for $U \in F[n]$ then return

$$g_k = \frac{U(n)}{V(n)} f_k, \text{ otherwise return "no } qm\text{-hypergeometric solution of (5.1)}$$

exists". □

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